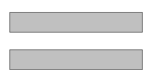
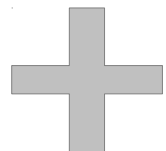




GK SOFTWARE



Java 8 ready

JDT Embraces Java 8

Stephan Herrmann

Committer:

- JDT
- Object Teams

Eclipse Track @ JavaLand 2014, Brühl, Germany

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THE RETAIL APPLICATION Company

- POS
- Store Integration
- Mobile Devices
- ...

...

Domain Specific Languages
Java (Server & Client)
Hardware Integration





- Andy Clement
- Steve Francisco
- Michael Rennie
- Olivier Thomann
- Curtis Windatt



- Walter Harley
- David Williams



- Jesper S. Møller



- Stephan Herrmann



- Dani Megert
- Markus Keller



- Jay Arthanareeswaran
- Anirban Chakarborty
- Manoj Palat
- Shankha Banerjee
- Manju Mathew
- Noopur Gupta
- Deepak Azad
- **Srikanth Sankaran**



<http://download.eclipse.org/eclipse/downloads>

http://wiki.eclipse.org/JDT/Eclipse_Java_8_Support_For_Kepler

- Java 8 features supported:
 - JSR 308 - Type Annotations.
 - JEP 120 - Repeating Annotations.
 - JEP 118 - Method Parameter Reflection.
 - JSR 269 - Pluggable Annotation Processor API & javax.lang.model API enhancements for Java 8.
 - JSR 335 – Lambda Expressions
 - Lambda expressions & method/constructor references
 - Support for “code carrying” interface methods
 - Enhanced target typing
new overload resolution & type inference



The Engine behind Lambda Expressions

- **Eclipse uses the Jikes Parser Generator**
 - requires an LALR(1) grammar
 - generates an automaton
- **Java 8 grammar**
 - is not LALR(1)
 - rewrite estimated to 10-15 person years
- **Rewrite could be avoided**
 - a number of advanced tricks
 - unorthodox communication between scanner & parser



- What we do in Java 7 (instance of an anonymous class):

```
new BinaryOperator<Integer>() {  
    @Override  
    public Integer apply(Integer i1, Integer i2) {  
        return i1+i2;  
    }  
}
```

- What want in Java 8 (lambda):

```
(i1, i2) -> i1+i2
```

- What can be omitted:

- the functional interface
- its generic type arguments
- the function signature
- ...

- How come this works?

- Lambdas are typically arguments to generic library functions

```
public static <T, K, U, M extends Map<K, U>>
Collector<T, ?, M> toMap(Function<? super T, ? extends K> keyMapper,
                        Function<? super T, ? extends U> valueMapper,
                        BinaryOperator<U> mergeFunction)
```

- Would you prefer to call:

```
Collectors.<Person,String,Integer,Map<String,Integer>>toMap(..)
```

OR: `Collectors.toMap(..)`

?

- Again: omitting “unnecessary” boiler plate
- Similar for instantiation of generic class (Java 7)

```
Collector<String,Integer> coll = new MyCollector<>();
```

- How *does* it work?

- Present since Java 5 – but completely rewritten
- Let the compiler collect *all available* type information
- Let it deduce the types
 - instantiation of type parameters of
 - generic methods
 - diamond instantiations
 - functional interface
 - lambda parameters
- If inference finds a solution, it is type safe – by definition
- If it doesn't find a solution ... ? ...

- Solving this little example

```
Map<String, Integer> test3(Stream<Person> persons) {  
    return persons.collect(Collectors.toMap(  
        p -> p.getLastName(),  
        p -> p.getAge(),  
        (i1, i2) -> i1+i2));  
}
```

- Produces these constraints:



- Solving th

```
Map<Strin
return
}
}
```

- Produces

TypeBound K#3 := java.lang.String
Dependency K#3 = K#3
TypeBound K#3 = java.lang.String
TypeBound K#3 <: java.lang.Object
Dependency R#0 = java.util.Map<K#3,U#4>
Dependency R#0 = java.util.Map<java.lang.String,U#4>
Dependency R#0 = java.util.Map<K#3,java.lang.Integer>
TypeBound R#0 = java.util.Map<java.lang.String,java.lang.Integer>
TypeBound R#0 <: java.util.Map<java.lang.String,java.lang.Integer>
TypeBound R#0 <: java.lang.Object
TypeBound T#2 := Person
Dependency T#2 = T#2
TypeBound T#2 = Person
TypeBound T#2 <: java.lang.Object
Dependency Map<K#3,U#4>#6 = java.util.Map<K#3,U#4>
Dependency Map<K#3,U#4>#6 = java.util.Map<java.lang.String,U#4>
Dependency Map<K#3,U#4>#6 = java.util.Map<K#3,java.lang.Integer>
TypeBound Map<K#3,U#4>#6 = java.util.Map<java.lang.String,java.lang.Integer>
Dependency Map<K#3,U#4>#6 = R#0
TypeBound Map<K#3,U#4>#6 <: java.lang.Object
TypeBound Map<K#3,U#4>#6 <: java.util.Map<java.lang.String,java.lang.Integer>
TypeBound A#1 = java.lang.Object
TypeBound A#1 <: java.lang.Object
TypeBound U#4 := java.lang.Integer
TypeBound U#4 = java.lang.Integer
Dependency U#4 = U#4
TypeBound U#4 <: java.lang.Object
TypeBound ?#5 = java.lang.Object
Dependency ?#5 = A#1
TypeBound ?#5 <: java.lang.Object

- Information flows:
 - Used for inference:
 - receiver type
 - argument types
 - target type



```
Target var = receiver.method(argument)
```

- Information flows:
 - Used for inference:
 - receiver type
 - argument types
 - target type
 - Never
 - what's right of the next dot

```
Target var = receiver.method(argument) .chain(a2);
```



- Information flows:

- Used for inference:

- receiver type
 - argument types
 - target type

- Never

- what's right of the next dot

- Nested inference

- expressions nested as call arguments
 - combine inner and outer to one inference

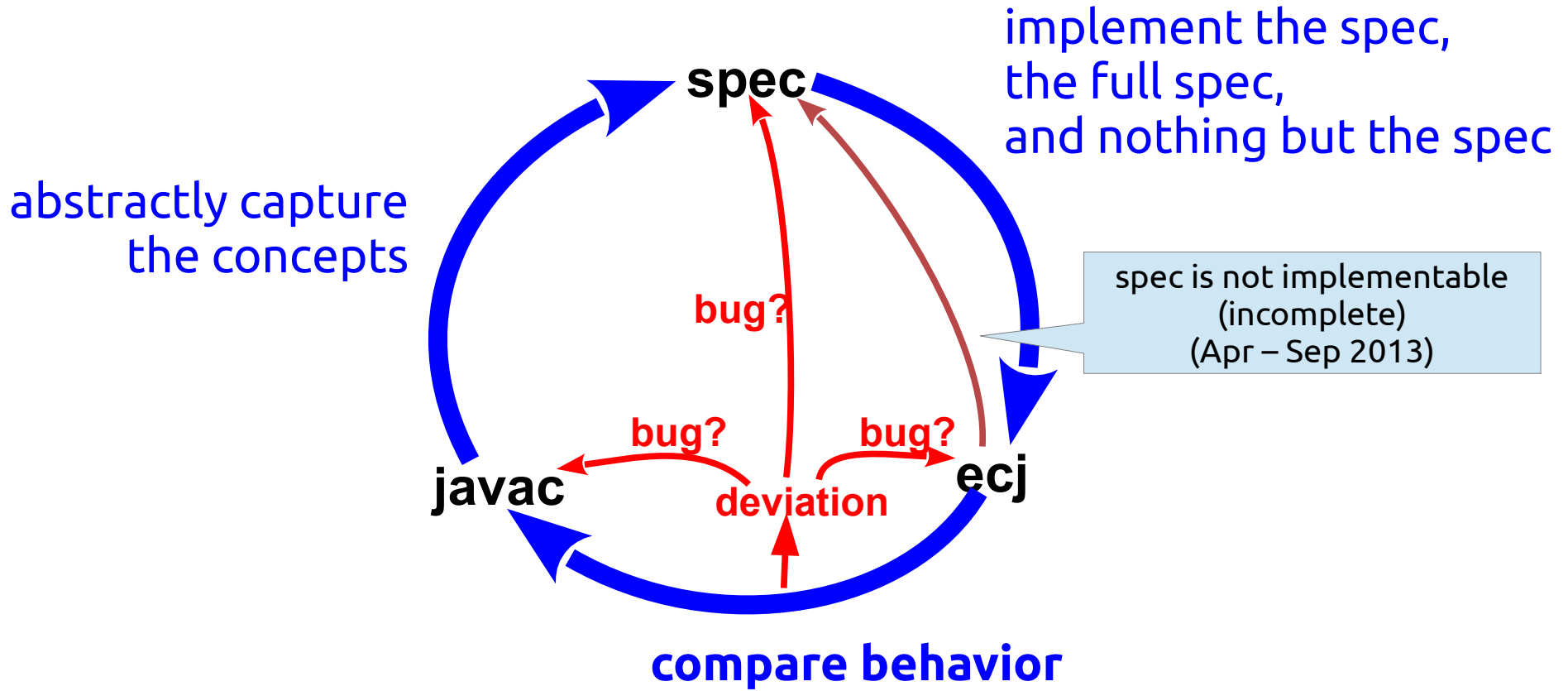
```
Map<String, Integer> test3(Stream<Person> ps) {  
    return ps  
        .collect(  
            Collectors.toMap(  
                p -> p.getLastName(),  
                p -> p.getAge(),  
                (i1, i2) -> i1+i2));  
}
```

If it fails ... ?

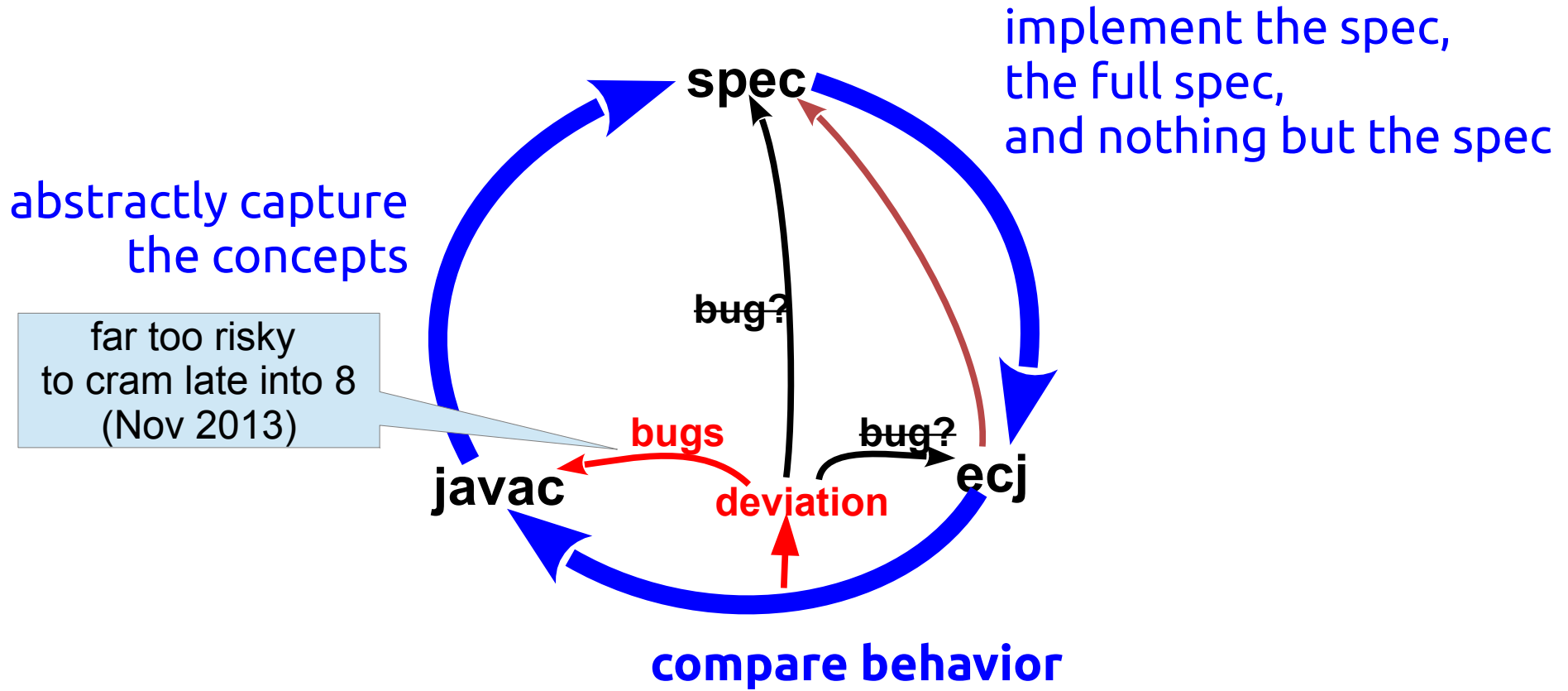
- To make sure you know what you're doing:
 - insert explicit types
- There's no guarantee that inference finds every solution
 - *“reduction is not completeness-preserving” (JLS 18.2.)*
- Inference is constraint solving
 - unlike common sense reasoning
 - if it fails, it's **hard** to say, why
 - ✘ The method foo(T1) in the type C is not applicable for the arguments (..)
 - ✘ Type mismatch: cannot convert from Foo to Bar
 - Better error reporting: future work
- Simplification for users
 - bought by tremendous complexity in the compiler



The Process behind a new Java Feature




Java Quality Assurance



Raw Types in Type Inference

- Where are unchecked conversions allowed?
 - List <: List<String> ?
 - spec: NO
 - javac: maybe, sometimes
 - ecj ?
 - raw types
 - supported in Java 5 to allow migration
 - some programs using raw types will eventually be rejected
- Serving two masters – JLS & javac?
 - should we exactly copy javac behavior?
 - is precisely implementing the spec better?
 - competition helps improve quality

- JLS 8 is more precise than previous versions
 - will be a good judge in case of doubt
- Ecj implementation is very close to the spec
 - bugs can be easily identified and fixed
 - 1 unimplemented branch, if you see:

you win! :)



- DEMO JDT Support for Lambda
 - Quick assist
 - to / from lambda
 - expression ↔ block
 - Search / Type hierarchy / Refactoring
 - are lambda aware
 - Hover
 - to see the functional type



Breathing Life into Types

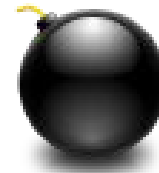
- **Java 5: annotate declarations**
 - ElementType: packages, classes, fields, methods, locals ...
- **Java 8: annotate types**
 - ElementType.TYPE_USE
 - ElementType.TYPE_PARAMETER



So what?

- Dynamically typed languages

- anything goes
- but may fail at runtime
- e.g.: “method not understood”



```
dog1 = new Dog();
dog1.bark();
```

```
dog2 = new Object();
dog2.bark();
```

- Type = Constraints on values

To *statically* detect anomalies

- missing capability
- incompatible assignment
- undeclared capability



```
Dog dog1 = new Dog();
dog1.bark();
```

```
Object dog2 = new Object();
dog2.bark();
```

```
Dog dog3 = new Object();
dog3.bark();
```

```
Object dog4 = new Dog();
dog4.bark();
```


Why Care About Types?

- **Dynamic**

- anything
- but m
- e.g.: “

Annotate source code:

- make assumptions explicit
- convince the compiler that code is safe

- **Type = Constraints on values**

To *statically* detect anomalies

- missing capability
- incompatible assignment
- undeclared capability

```
Dog dog1 = new Dog();  
dog1.bark();
```

```
Object dog2 = new Object();  
dog2.bark();
```

```
Dog dog3 = new Object();  
dog3.bark();
```

```
Object dog4 = new Dog();  
dog4.bark();
```

- Constraint checking avoids errors

- No Such Method / Field

- Basic statically typed OO

- ClassCastException

- Generics

- ??Exception

- SWTException("Invalid thread access")
- ...
- **NullPointerException**

```
void letemBark(Vector dogs) {  
    Dog aDog = (Dog) dogs.get(0);  
    aDog.bark();  
}
```

- Make it easier to add new constraints
 - Only one new syntax for all kinds of constraints
- Make it easier to add new type checkers
 - Checker Framework (Michael Ernst – U of Washington)
- Examples
 - @NonNull
 - @Interned equals(== , equals)
 - @Immutable value cannot change (Java 5 ?)
 - @ReadOnly value cannot change *via this reference*
 - @UI code requires to run on the UI thread

Can't Java 5 Do All This?

```
@Target(ElementType.PARAMETER)
@interface NonNull5 {}

void java5(@NonNull5 String arg);
```

arg is qualified to be non-null

```
@Target(ElementType.METHOD)
@interface NonNull5 {}

@NonNull5 String java5();
```

java5 is qualified to be non-null

```
@Target(ElementType.TYPE_USE)
@interface NonNull8 {}

void java8(@NonNull8 String arg);
```

String is qualified to be non-null

```
@Target(ElementType.TYPE_USE)
@interface NonNull8 {}

@NonNull8 String java8();
```

String is qualified to be non-null

- We've been lying about the method result
 - but we can't lie about everything, e.g.:

```
void printFirstDog(@NonNull List<Dog> dogs) {
    dogs.get(0).bark();
}
```

NPE?


```
void good() {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add("Hello");  
    for (String elem : l1)  
        System.out.println(elem.toUpperCase());  
  
    @NonNull List<@Nullable String> l2 = new ArrayList<>();  
    l2.add(null);  
    for (String unknown : l2)  
        if (unknown != null)  
            System.out.println(unknown.toUpperCase());  
}
```

l1 cannot contain
null elements

l2 can contain
null elements

```
void bad(List<String> unknown, List<@Nullable String> withNulls) {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add(null);  
    String  
    if (fir  
        l1 = unknown;  
        l1 = withNulls;  
  
    String canNull = withNulls.get(0);  
    System.out.println(canNull.toUpperCase());  
}
```

✘ Null type mismatch: required '@NonNull String' but the provided value is null

```
void bad(List<String> unknown, List<@Nullable String> withNulls) {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add(null);  
    String first = l1.get(0);  
    if (first == null) return;  
  
    l1 =  Null comparison always yields false:  
    l1 = The variable first cannot be null at this location  
  
    String canNull = withNulls.get(0);  
    System.out.println(canNull.toUpperCase());  
}
```


```
void bad(List<String> unknown, List<@Nullable String> withNulls) {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add(null);  
    String first = l1.get(0);  
    if (first == null) return;  
  
    l1 = unknown;  
    l1 =  
    String  
    System.out.println(canNull.toUpperCase());  
}
```

🔒 Null type safety (type annotations): The expression of type 'List<String>' needs unchecked conversion to conform to '@NonNull List<@NonNull String>'


```
void bad(List<String> unknown, List<@Nullable String> withNulls) {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add(null);  
    String first = l1.get(0);  
    if (first == null) return;  
  
    l1 = unknown;  
    l1 = withNulls;  
  
    String  
    System  
}
```

❌ Null type mismatch (type annotations):
required '@NonNull List<@NonNull String>'
but this expression has type 'List<@Nullable String>'

```
void bad(List<String> unknown, List<@Nullable String> withNulls) {  
    @NonNull List<@NonNull String> l1 = new ArrayList<>();  
    l1.add(null);  
    String first = l1.get(0);  
    if (first == null) return;  
  
    l1 = unknown;  
    l1 = withNulls;  
  
    String canNull = withNulls.get(0);  
    System.out.println(canNull.toUpperCase(););  
}
```

 Potential null pointer access:
The variable canNull may be null at this location

- **Null Annotations**
 - org.eclipse.jdt.annotation_2.0.0
 - @NonNull, @Nullable specify Target(TYPE_USE)
 - @NonNullByDefault: more fine tuning
- **Null Analysis (per compiler option)**
 - Nullness is an integral part of the type system
 - Fine tuned defaults only in Luna
 - TODO: Strict checking against type variables

- Null Annotations
- Null Analysis (per compiler option)
- Planned: @Uninterned
 - Detect accidental comparison using == or !=
- Proposed: @UiEffect, @Ui ...
 - by [Colin S. Gordon, Werner Dietl, Michael D. Ernst, and Dan Grossman]
 - SWT: bye, bye, "Invalid thread access"

- Null Annotations
- Null Analysis (per compiler option)
- Planned: @Uninterned
- Proposed: @UiEffect, @Ui ...
- JDT/UI
 - OK: completion, refactoring, etc.
 - TODO: show in hover
 - TODO: update / add more quickfixes

- Null Annotations
- Null Analysis (per compiler option)
- Planned: @Uninterned
- Proposed: @UiEffect, @Ui ...
- JDT/UI

Do You care about Types?



Breathing Life into Objects (beyond Java)

- Algorithms

- lambda expressions
- library updates for lambdas

λ

- Types

- type annotations

~~NPE~~

- Objects? – Connecting objects to a module / a system?

- references
 - bad for analyzeability / modularity
- inheritance
 - doesn't scale
 - rigid

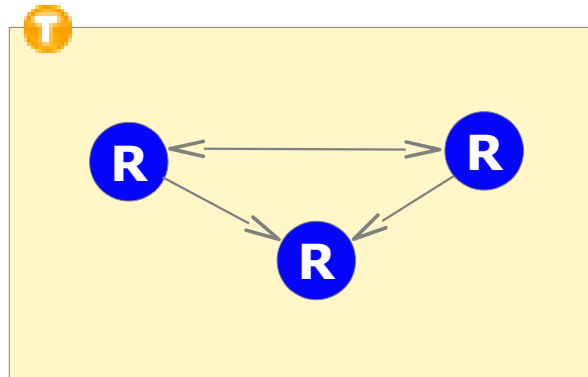
- 2002: Research on Language Design beyond standard OO
 - Object Teams – programming with roles and teams
 - OT/J – extends Java



- 2010: Eclipse Object Teams Project established
 - Outgrown the research niche
 - Mature language (OT/J) & IDE (OTDT)
 - Object Teams Development Tooling written in OT/J

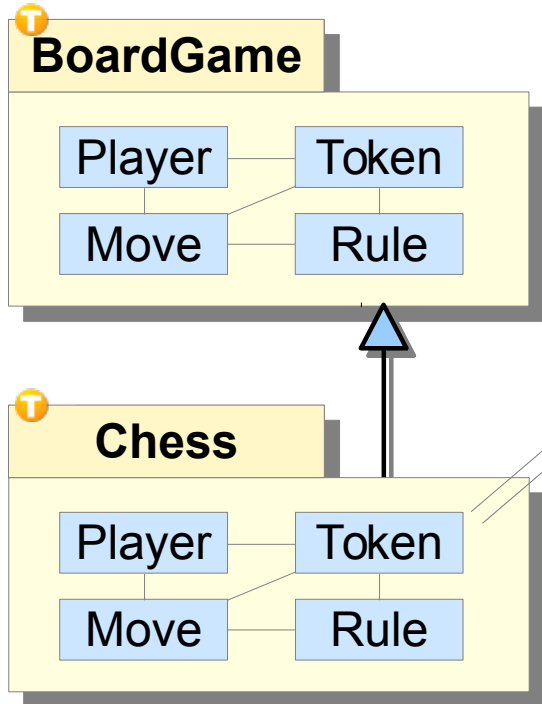


- OO: Message passing among objects



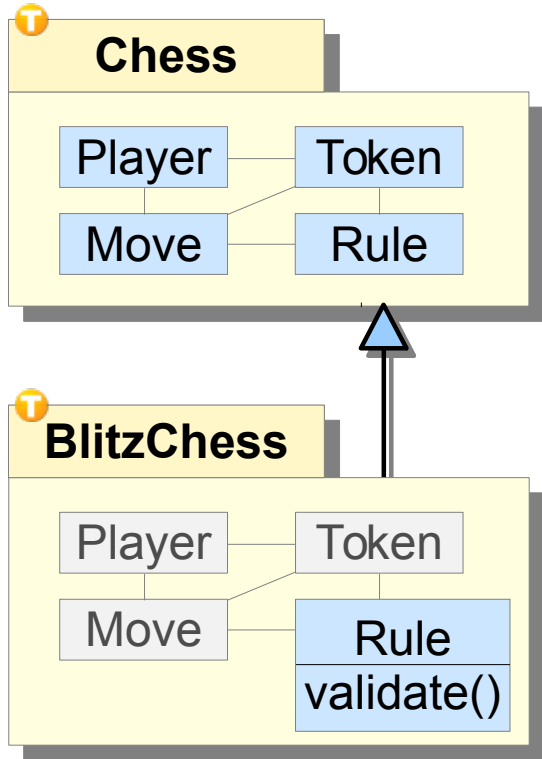
- Team: A group of interacting Roles
 - A team instance **owns** its role instances
 - Team provides a boundary for **isolation**

```
team class Company {
    class Employee {
    }
}
```



impossible to mix elements from different games

- Members of a team are **virtual classes**
 - consistent refinement of all members



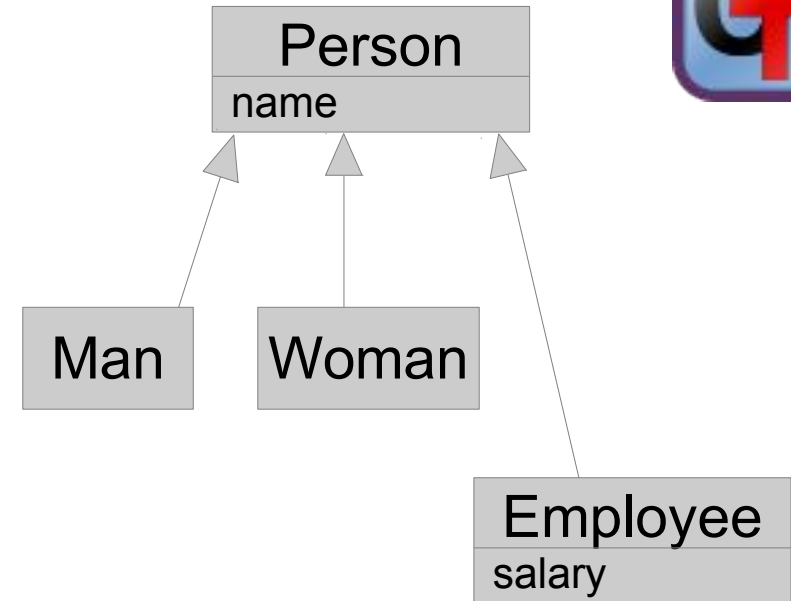
- consistent refinement of all members
- deep overriding
- flexible & modular



- **Inheritance** is great, but ...

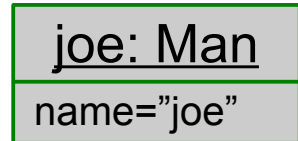
A text book example:

- A man/woman **is a** person, OK
- An employee **is a** person, OK?
 - Male/female Employees?
 - Born as an employee?
 - Dying when loosing the job?
 - Several jobs, yet only one salary?





- playedBy relationship

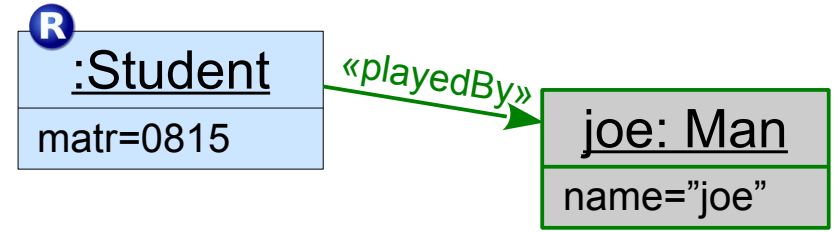




- playedBy relationship



- Advantages:
 - Dynamism:**
roles can come and go
(same base object)



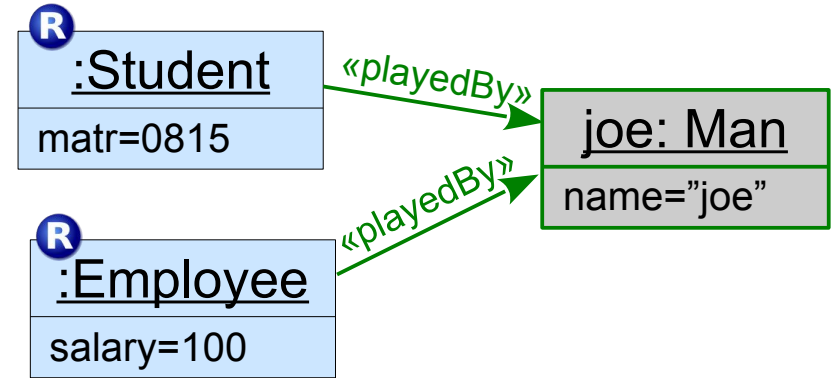


- playedBy relationship



- Advantages:

- **Dynamism:**
roles can come and go
(same base object)
- **Multiplicities:**
one base can play several roles
(different/same role types)



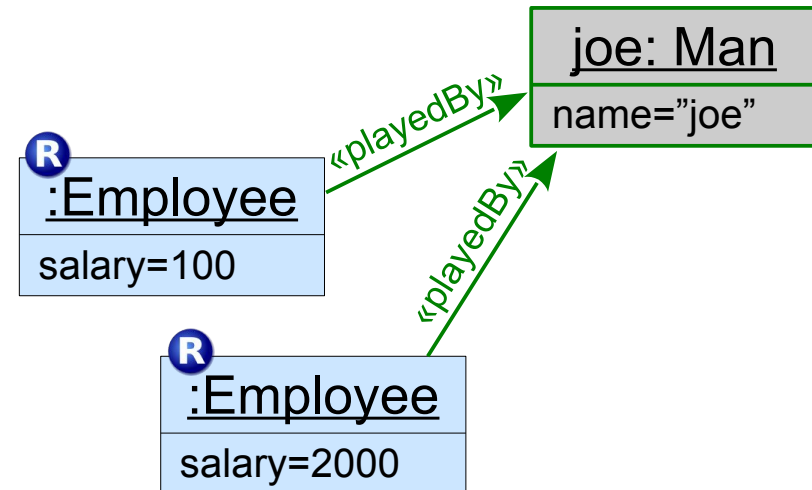


- playedBy relationship



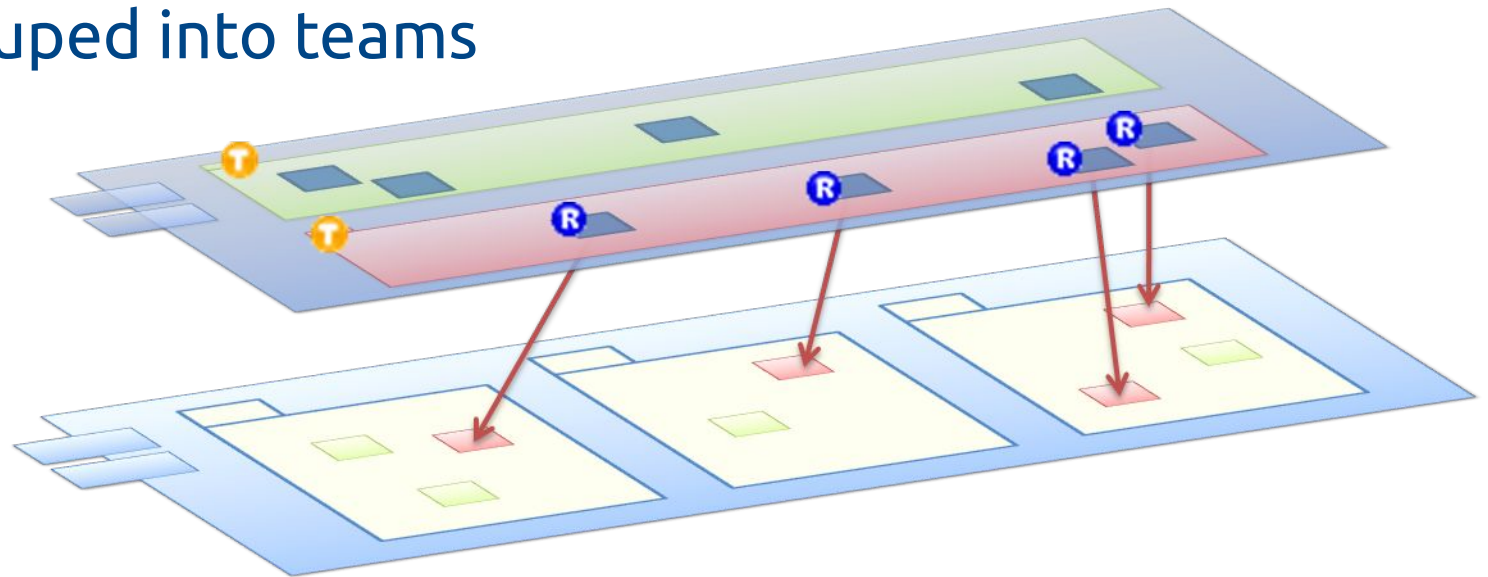
- Advantages:

- **Dynamism:**
roles can come and go
(same base object)
- **Multiplicities:**
one base can play several roles
(different/same role types)





- Use role playing for
 - dynamic specialization
 - modes / behavior that should be turn on/off at runtime
 - unanticipated adaptation
- Roles are grouped into teams



- **Methods & algorithms**

- lambda expressions
- library updates for lambdas

- **Types**

- type annotations

- **Connecting objects to a module / a system**

- role containment
- team inheritance
- role playing

